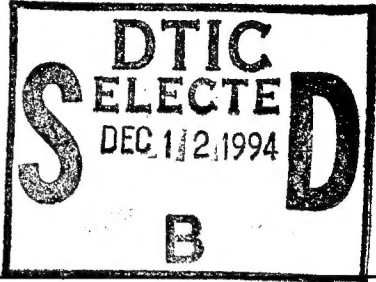# REPORT DOCUMENTATION PAGE

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1. AGENCY USE (Leave	2. REPORT	3. REPORT TYPE AND DATES

4. TITLE AND: UNISYS Corporation
Compiler: IntergrAda for Windows NT, Version 1.0
Host/Target: Intel Deskside Server with Intel 60 MHz Pentium cpu (under Microsoft Windows NT Server, Version 3.5)

5. FUNDING

DTIC
SELECTED
DEC 12,1994
B

6. Authors: The National Institute of Standards and Technology

7. PERFORMING ORGANIZATION NAME (S) AND:
Computer Systems Laboratory (CSL)
National Institute of Standards and Technology
Building 225, Room A266
Gaithersburg, MD 20899

8. PERFORMING ORGANIZATION

9. SPONSORING/MONITORING AGENCY NAME(S) AND:
Ada Joint Program Office, Defense Information Systems Agency
Code TXEA, 701 S. Courthouse Rd.
Arlington, VA 22204-2199

10. SPONSORING/MONITORING AGENCY

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VCL#: 940902S1.11376, AVF#: 94uni501_1

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17 SECURITY CLASSIFICATION	18. SECURITY	19. SECURITY CLASSIFICATION	20. LIMITATION OF
UNCLASSIFIED	**UNCLASSIFIED**	**UNCLASSIFIED**	**UNCLASSIFIED**

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AVF Control Number: NIST94UNI501_1_1.11
DATE COMPLETED
    BEFORE ON-SITE: 94-08-31
    AFTER   ON-SITE: 94-09-06
    REVISIONS: 94-09-14

Ada   COMPILER
VALIDATION SUMMARY REPORT:
Certificate Number: 940902S1.11376
UNISYS Corporation
IntegrAda for Windows NT, Version 1.0
Intel Deskside Server with Intel Pentium 60 MHz =>
Intel Deskside Server with Intel Pentium 60 MHz

Prepared By:
Software Standards Validation Group
Computer Systems Laboratory
National Institute of Standards and Technology
Building 225, Room A266
Gaithersburg, Maryland 20899
U.S.A.

19941202 190

Certificate Information

The following Ada implementation was tested and determined to pass ACVC 1.11. Testing was completed on September 2, 1994.

    Compiler Name and Version: IntegrAda for Windows NT, Version 1.0

    Host Computer System: Intel Deskside Server with Intel Pentium
                                  60 MHz under Microsoft Windows NT Server,
                                  Version 3.5

    Target Computer System: Intel Deskside Server with Intel Pentium
                                  60 MHz under Microsoft Windows NT Server,
                                  Version 3.5

See section 3.1 for any additional information about the testing environment.

As a result of this validation effort, Validation Certificate 940902S1.11376 is awarded to UNISYS Corporation. This certificate expires 2 years after ANSI/MIL-STD-1815B is approved by ANSI.

This report has been reviewed and is approved.

_____   _____
Ada Validation Facility   Ada Validation Facility
Dr. David K. Jefferson   Mr. L. Arnold Johnson
Chief, Information Systems   Manager, Software Standards
    Engineering Division (ISED)   Validation Group
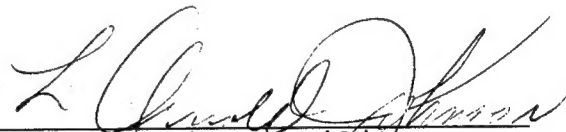   Computer Systems Laboratory (CSL)
            National Institute of Standards and Technology
                    Building 225, Room A266
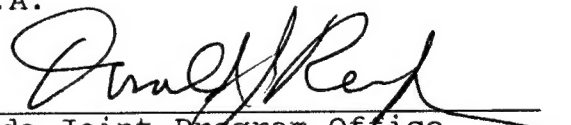   Gaithersburg, Maryland 20899
                        U.S.A.

_____   _____
Ada Validation Organization   Ada Joint Program Office
Director, Computer & Software   Donald J. Reifer
    Engineering Division   Director, Ada Joint Program Office
Institute for Defense Analyses   Defense Information Systems Agency,
Alexandria VA 22311   Center for Information Management
                                         Washington DC 20301
                        U.S.A.

DECLARATION OF CONFORMANCE

The following declaration of conformance was supplied by the customer.

Customer: UNISYS Corporation

Certificate Awardee: UNISYS Corporation

Ada Validation Facility: National Institute of Standards and
                                    Technology
                                Computer Systems Laboratory (CSL)
                                Software Standards Validation Group
                                Building 225, Room A266
                                Gaithersburg, Maryland 20899
                                U.S.A.

ACVC Version: 1.11

Ada Implementation:

   Compiler Name and Version: IntegrAda for Windows NT, Version 1.0

   Host Computer System: Intel Deskside Server with Intel Pentium
                                60 MHz under Microsoft Windows NT Server,
                                Version 3.5

   Target Computer System: Intel Deskside Server with Intel Pentium
                                60 MHz under Microsoft Windows NT Server,
                                Version 3.5

Declaration:

I the undersigned, declare that I have no knowledge of deliberate
deviations from the Ada Language Standard ANSI/MIL-STD-1815A ISO
8652-1987 in the implementation listed above.

_____   _9/7/54_____
Customer Signature   Date
Company UNISYS Corporation
Title Manager, Contracts and Pricing

_____   _5/7/54_____
Certificate Awardee Signature   Date
Company UNISYS Corporation
Title Manager, Contracts and Pricing

Accession For	
NTIS GRA&I	☑
DTIC TAB	☐
Unannounced	☐
Justification	
By_____	
Distribution/	
Availability Codes	
Dist	Avail and/or Special
A-1	

TABLE OF CONTENTS

CHAPTER 1

INTRODUCTION

The Ada implementation described above was tested according to the Ada Validation Procedures [Pro92] against the Ada Standard [Ada83] using the current Ada Compiler Validation Capability (ACVC). This Validation Summary Report (VSR) gives an account of the testing of this Ada implementation. For any technical terms used in this report, the reader is referred to [Pro92]. A detailed description of the ACVC may be found in the current ACVC User's Guide [UG89].

1.1 USE OF THIS VALIDATION SUMMARY REPORT

Consistent with the national laws of the originating country, the Ada Certification Body may make full and free public disclosure of this report. In the United States, this is provided in accordance with the "Freedom of Information Act" (5 U.S.C. #552). The results of this validation apply only to the computers, operating systems, and compiler versions identified in this report.

The organizations represented on the signature page of this report do not represent or warrant that all statements set forth in this report are accurate and complete, or that the subject implementation has no nonconformities to the Ada Standard other than those presented. Copies of this report are available to the public from the AVF which performed this validation or from:

> National Technical Information Service
> 5285 Port Royal Road
> Springfield, Virginia 22161
> U.S.A.

Questions regarding this report or the validation test results should be directed to the AVF which performed this validation or to:

> Ada Validation Organization
> Computer and Software Engineering Division
> Institute for Defense Analyses
> 1801 North Beauregard Street
> Alexandria, Virginia 22311-1772
> U.S.A.

1.2 REFERENCES

[Ada83] <u>Reference Manual for the Ada Programming Language</u>, ANSI/MIL-STD-1815A, February 1983 and ISO 8652-1987.

[Pro92] <u>Ada Compiler Validation Procedures</u>, Version 3.1, Ada Joint Program Office, August 1992.

[UG89] <u>Ada Compiler Validation Capability User's Guide</u>, 21 June 1989.

1.3 ACVC TEST CLASSES

Compliance of Ada implementations is tested by means of the ACVC. The ACVC contains a collection of test programs structured into six test classes: A, B, C, D, E, and L. The first letter of a test name identifies the class to which it belongs. Class A, C, D, and E tests are executable. Class B and class L tests are expected to produce errors at compile time and link time, respectively.

The executable tests are written in a self-checking manner and produce a PASSED, FAILED, or NOT APPLICABLE message indicating the result when they are executed. Three Ada library units, the packages REPORT and SPPRT13, and the procedure CHECK_FILE are used for this purpose. The package REPORT also provides a set of identity functions used to defeat some compiler optimizations allowed by the Ada Standard that would circumvent a test objective. The package SPPRT13 is used by many tests for Chapter 13 of the Ada Standard. The procedure CHECK_FILE is used to check the contents of text files written by some of the Class C tests for Chapter 14 of the Ada Standard. The operation of REPORT and CHECK_FILE is checked by a set of executable tests. If these units are not operating correctly, validation testing is discontinued.

Class B tests check that a compiler detects illegal language usage. Class B tests are not executable. Each test in this class is compiled and the resulting compilation listing is examined to verify that all violations of the Ada Standard are detected. Some of the class B tests contain legal Ada code which must not be flagged illegal by the compiler. This behavior is also verified.

Class L tests check that an Ada implementation correctly detects violation of the Ada Standard involving multiple, separately compiled units. Errors are expected at link time, and execution is attempted.

In some tests of the ACVC, certain macro strings have to be replaced by implementation-specific values--for example, the

largest integer. A list of the values used for this implementation is provided in Appendix A. In addition to these anticipated test modifications, additional changes may be required to remove unforeseen conflicts between the tests and implementation-dependent characteristics. The modifications required for this implementation are described in section 2.3.

For each Ada implementation, a customized test suite is produced by the AVF. This customization consists of making the modifications described in the preceding paragraph, removing withdrawn tests (see section 2.1) and, possibly some inapplicable tests (see Section 3.2 and [UG89]).

In order to pass an ACVC an Ada implementation must process each test of the customized test suite according to the Ada Standard.

1.4 DEFINITION OF TERMS

Ada Compiler   The software and any needed hardware that
                            have to be added to a given host and target
                            computer system to allow transformation of
                            Ada programs into executable form and
                            execution thereof.

Ada Compiler   The means for testing compliance of Ada
Validation   implementations, Validation consisting of
Capability (ACVC)   the test suite, the support programs, the
                            ACVC Capability User's Guide and the
                            template for the validation summary (ACVC)
                            report.

Ada Implementation   An Ada compiler with its host computer
                            system and its target computer system.

Ada Joint Program   The part of the certification body which
Office (AJPO)   provides policy and guidance for the Ada
                            certification Office system.

Ada Validation   The part of the certification body which
Facility (AVF)   carries out the procedures required to
                            establish the compliance of an Ada
                            implementation.

Ada Validation   The part of the certification body that
Organization (AVO)   provides technical guidance for operations
                            of the Ada certification system.

Compliance of an   The ability of the implementation to pass an
Ada Implementation   ACVC version.

Computer System	A functional unit, consisting of one or more computers and associated software, that uses common storage for all or part of a program and also for all or part of the data necessary for the execution of the program; executes user- written or user-designated programs; performs user-designated data manipulation, including arithmetic operations and logic operations; and that can execute programs that modify themselves during execution. A computer system may be a stand-alone unit or may consist of several inter-connected units.
Conformity	Fulfillment by a product, process, or service of all requirements specified.
Customer	An individual or corporate entity who enters into an agreement with an AVF which specifies the terms and conditions for AVF services (of any kind) to be performed.
Declaration of Conformance	A formal statement from a customer assuring that conformity is realized or attainable on the Ada implementation for which validation status is realized.
Host Computer System	A computer system where Ada source programs are transformed into executable form.
Inapplicable Test	A test that contains one or more test objectives found to be irrelevant for the given Ada implementation.
ISO	International Organization for Standardization.
LRM	The Ada standard, or Language Reference Manual, published as ANSI/MIL-STD-1815A -1983 and ISO 8652-1987. Citations from the LRM take the form "<section>.<subsection>: <paragraph>."
Operating System	Software that controls the execution of programs and that provides services such as resource allocation, scheduling, input/output control, and data management. Usually, operating systems are predominantly software, but partial or complete hardware implementations are possible.
Target Computer System	A computer system where the executable form of Ada programs are executed.

Validated Ada Compiler	The compiler of a validated Ada implementation.
Validated Ada Implementation	An Ada implementation that has been validated successfully either by AVF testing or by registration [Pro92].
Validation	The process of checking the conformity of an Ada compiler to the Ada programming language and of issuing a certificate for this implementation.
Withdrawn Test	A test found to be incorrect and not used in conformity testing. A test may be incorrect because it has an invalid test objective, fails to meet its test objective, or contains erroneous or illegal use of the Ada programming language.

CHAPTER 2

IMPLEMENTATION DEPENDENCIES

2.1 WITHDRAWN TESTS

Some tests are withdrawn by the AVO from the ACVC because they do
not conform to the Ada Standard. The following 104 tests had been
withdrawn by the Ada Validation Organization (AVO) at the time of
validation testing. The rationale for withdrawing each test is
available from either the AVO or the AVF. The publication date for
this list of withdrawn tests is 93-11-22.

B27005A	E28005C	B28006C	C32203A	C34006D	C35507K
C35507L	C35507N	C35507O	C35507P	C35508I	C35508J
C35508M	C35508N	C35702A	C35702B	C37310A	B41308B
C43004A	C45114A	C45346A	C45612A	C45612B	C45612C
C45651A	C46022A	B49008A	B49008B	A54B02A	C55B06A
A74006A	C74308A	B83022B	B83022H	B83025B	B83025D
B83026B	C83026A	C83041A	B85001L	C86001F	C94021A
C97116A	C98003B	BA2011A	CB7001A	CB7001B	CB7004A
CC1223A	BC1226A	CC1226B	BC3009B	BD1B02B	BD1B06A
AD1B08A	BD2A02A	CD2A21E	CD2A23E	CD2A32A	CD2A41A
CD2A41E	CD2A87A	CD2B15C	BD3006A	BD4008A	CD4022A
CD4022D	CD4024B	CD4024C	CD4024D	CD4031A	CD4051D
CD5111A	CD7004C	ED7005D	CD7005E	AD7006A	CD7006E
AD7201A	AD7201E	CD7204B	AD7206A	BD8002A	BD8004C
CD9005A	CD9005B	CDA201E	CE2107I	CE2117A	CE2117B
CE2119B	CE2205B	CE2405A	CE3111C	CE3116A	CE3118A
CE3411B	CE3412B	CE3607B	CE3607C	CE3607D	CE3812A
CE3814A	CE3902B				

2.2 INAPPLICABLE TESTS

A test is inapplicable if it contains test objectives which are
irrelevant for a given Ada implementation. The inapplicability
criteria for some tests are explained in documents issued by ISO
and the AJPO known as Ada Commentaries and commonly referenced in
the format AI-ddddd. For this implementation, the following tests
were determined to be inapplicable for the reasons indicated;
references to Ada Commentaries are included as appropriate.

The following 201 tests have floating-point type declarations
requiring more digits than SYSTEM.MAX_DIGITS:

> C24113L..Y (14 tests)   C35705L..Y (14 tests)
> C35706L..Y (14 tests)   C35707L..Y (14 tests)

```
         C35708L..Y (14 tests)      C35802L..Z (15 tests)
         C45241L..Y (14 tests)      C45321L..Y (14 tests)
         C45421L..Y (14 tests)      C45521L..Z (15 tests)
         C45524L..Z (15 tests)      C45621L..Z (15 tests)
         C45641L..Y (14 tests)      C46012L..Z (15 tests)
```

The following 20 tests check for the predefined type LONG_INTEGER;
for this implementation, there is no such type:

```
    C35404C      C45231C      C45304C      C45411C      C45412C
    C45502C      C45503C      C45504C      C45504F      C45611C
    C45613C      C45614C      C45631C      C45632C      B52004D
    C55B07A      B55B09C      B86001W      C86006C      CD7101F
```

C35713B, C45423B, B86001T, and C86006H check for the predefined
type SHORT_FLOAT; for this implementation, there is no such type.

C35713D and B86001Z check for a predefined floating-point type with
a name other than FLOAT, LONG_FLOAT, or SHORT_FLOAT; for this
implementation, there is no such type.

C45531M..P and C45532M..P (8 tests) check fixed-point operations
for types that require a SYSTEM.MAX_MANTISSA of 47 or greater; for
this implementation, MAX_MANTISSA is less than 47.

C45536A, C46013B, C46031B, C46033B, and C46034B contain length
clauses that specify values for 'SMALL that are not powers of two
or ten; this implementation does not support such values for
'SMALL.

C45624A..B (2 tests) check that the proper exception is raised if
MACHINE_OVERFLOWS is FALSE for floating point types and the results
of various floating-point operations lie outside the range of the
base type; for this implementation, MACHINE_OVERFLOWS is TRUE.

B86001Y uses the name of a predefined fixed-point type other than
type DURATION; for this implementation, there is no such type.

C96005B uses values of type DURATION's base type that are outside
the range of type DURATION; for this implementation, the ranges are
the same.

CD1009C checks whether a length clause can specify a non-default
size for a floating-point type; this implementation does not
support such sizes.

CD2A53A checks operations of a fixed-point type for which a length
clause specifies a power-of-ten TYPE'SMALL; this implementation
does not support decimal 'SMALLs. (See section 2.3.)

CD2A84A, CD2A84E, CD2A84I..J (2 tests), and CD2A84O use length

clauses to specify non-default sizes for access types; this implementation does not support such sizes.

BD8001A, BD8003A, BD8004A..B (2 tests), and AD8011A use machine code insertions; this implementation provides no package MACHINE_CODE.

The 18 tests listed in the following table check that USE_ERROR is raised if the given file operations are not supported for the given combination of mode and access method; this implementation supports these operations.

Test	File Operation	Mode	File Access Method
CE2102E	CREATE	OUT_FILE	SEQUENTIAL_IO
CE2102F	CREATE	INOUT_FILE	DIRECT_IO
CE2102J	CREATE	OUT_FILE	DIRECT_IO
CE2102N	OPEN	IN_FILE	SEQUENTIAL_IO
CE2102O	RESET	IN_FILE-.	SEQUENTIAL_IO
CE2102P	OPEN	OUT_FILE	SEQUENTIAL_IO
CE2102Q	RESET	OUT_FILE	SEQUENTIAL_IO
CE2102R	OPEN	INOUT_FILE	DIRECT_IO
CE2102S	RESET	INOUT_FILE	DIRECT_IO
CE2102T	OPEN	IN_FILE	DIRECT_IO
CE2102U	RESET	IN_FILE	DIRECT_IO
CE2102V	OPEN	OUT_FILE	DIRECT_IO
CE2102W	RESET	OUT_FILE	DIRECT_IO
CE3102F	RESET	Any Mode	TEXT_IO
CE3102G	DELETE	--------	TEXT_IO
CE3102I	CREATE	OUT_FILE	TEXT_IO
CE3102J	OPEN	IN_FILE	TEXT_IO
CE3102K	OPEN	OUT_FILE	TEXT_IO

The 3 tests listed in the following table check the given file operations for the given combination of mode and access method; this implementation does not support these operations.

Test	File Operation	Mode	File Access Method
CE2105A	CREATE	IN_FILE	SEQUENTIAL_IO
CE2105B	CREATE	IN_FILE	DIRECT_IO
CE3109A	CREATE	IN_FILE	TEXT_IO

CE2203A checks that WRITE raises USE_ERROR if the capacity of an external sequential file is exceeded; this implementation cannot restrict file capacity.

EE2401D, and EE2401G use instantiations of DIRECT_IO with unconstrained array and record types; this implementation raises USE_ERROR on the attempt to create a file of such types.

CE2401H uses instantiations of DIRECT_IO with unconstrained array and record types; this implementation raises USE_ERROR on the attempt to create a file of such types.

CE2403A checks that WRITE raises USE_ERROR if the capacity of an external direct file is exceeded; this implementation cannot restrict file capacity.

CE3304A checks that SET_LINE_LENGTH and SET_PAGE_LENGTH raise USE_ERROR if they specify an inappropriate value for the external file; there are no inappropriate values for this implementation.

CE3413B checks that PAGE raises LAYOUT_ERROR when the value of the page number exceeds COUNT'LAST; for this implementation, the value of COUNT'LAST is greater than 150000, making the checking of this objective impractical.

.2.3 TEST MODIFICATIONS

Modifications (see section 1.3) were required for 19 tests.

The following tests were split into two or more tests because this implementation did not report the violations of the Ada Standard in the way expected by the original tests.

B23004A	B23007A	B23009A	B25002A	B26005A
B28003A	B32202A	B32202B	B32202C	B37004A
B61012A	B95069A	B95069B	BA1101B	BC2001D
BC3009A	BC3009C			

BA2001E was graded passed by Evaluation Modification as directed by the AVO. The test expects that duplicate names of subunits with a common ancestor will be detected as compilation errors; this implementation detects the errors at link time, and the AVO ruled that this behavior is acceptable.

CD2A53A was graded inapplicable by Evaluation Modification as directed by the AVO. The test contains a specification of a power-of-10 value as 'SMALL for a fixed-point type. The AVO ruled that, under ACVC 1.11, support of decimal 'SMALLs may be omitted.

CHAPTER 3

PROCESSING INFORMATION

3.1   TESTING ENVIRONMENT

The Ada implementation tested in this validation effort is described adequately by the information given in the initial pages of this report.

For technical information about this Ada implementation, contact:

> Mr. Michael Rowland
> 8008 Westpark Drive
> Mclean, VA 22102 (U.S.A.)

For sales information about this Ada implementation, contact:

> Mr. Thomas Breves
> 8008 Westpark Drive
> Mclean, VA 22102 (U.S.A.)

Testing of this Ada implementation was conducted at the customer's site by a validation team from the AVF.

3.2   SUMMARY OF TEST RESULTS

An Ada Implementation passes a given ACVC version if it processes each test of the customized test suite in accordance with the Ada Programming Language Standard, whether the test is applicable or inapplicable; otherwise, the Ada Implementation fails the ACVC [Pro92].

For all processed tests (inapplicable and applicable), a result was obtained that conforms to the Ada Programming Language Standard.

The list of items below gives the number of ACVC tests in various categories.   All tests were processed, except those that were withdrawn because of test errors (item b; see section 2.1), those that require a floating-point precision that exceeds the implementation's maximum precision (item e; see section 2.2), and those that depend on the support of a file system--if none is supported (item d). All tests passed, except those that are listed in sections 2.1 and 2.2 (counted in items b and f, below).

    a) Total Number of Applicable Tests   3782

    b) Total Number of Withdrawn Tests   104
    c) Processed Inapplicable Tests   284

d)	Non-Processed I/O Tests	0	
e)	Non-Processed Floating-Point Precision Tests	0	

f)	Total Number of Inapplicable Tests	284	(c+d+e)
g)	Total Number of Tests for ACVC 1.11	4170	(a+b+f)

3.3 TEST EXECUTION

A magnetic tape containing the customized test suite (see section
1.3) was taken on-site by the validation team for processing. The
contents of the magnetic tape were loaded directly onto the
host/target computer.

After the test files were loaded onto the host computer, the full
set of tests was processed by the Ada implementation.

The tests were compiled, linked, and executed on the host/target
computer system, as appropriate.

Testing was performed using command scripts provided by the
customer and reviewed by the validation team. See Appendix B for
a complete listing of the processing options for this
implementation. It also indicates the default options. The
options invoked explicitly for validation testing during this test
were:

CHECKS	=> ALL	Generate all execution checks.
GENERICS	=> STUBS	Do not inline generics.
TASKING	=> YES	Allow tasking.
MEMORY	=> 500	Amount of internal buffers shared by compile virtual memory.
STACK	=> 20480	Boundary size determining whether an dynamic object is allocated on the stack or in the map.
INLINE	=> PRAGMA	Inlining of subprograms by pragma INLINE.
REDUCTION	=> NONE	No optimization of check or loops.
EXPRESSIONS	=> NONE	No lowlevel optimization.

Test output, compiler and linker listings, and job logs were
captured on magnetic tape and archived at the AVF. The listings
examined on-site by the validation team were also archived.

APPENDIX A

MACRO PARAMETERS

This appendix contains the macro parameters used for customizing the ACVC. The meaning and purpose of these parameters are explained in [UG89]. The parameter values are presented in two tables. The first table lists the values that are defined in terms of the maximum input-line length, which is the value for $MAX_IN_LEN--also listed here. These values are expressed here as Ada string aggregates, where "V" represents the maximum input-line length.

Macro Parameter	Macro Value
$MAX_IN_LEN	<255> -- Value of V
$BIG_ID1	(1..V-1 => 'A', V => '1')
$BIG_ID2	(1..V-1 => 'A', V => '2')
$BIG_ID3	(1..V/2 => 'A') & '3' & (1..V-1-V/2 => 'A')
$BIG_ID4	(1..V/2 => 'A') & '4' & (1..V-1-V/2 => 'A')
$BIG_INT_LIT	(1..V-3 => '0') & "298"
$BIG_REAL_LIT	(1..V-5 => '0') & "690.0"
$BIG_STRING1	'"' & (1..V/2 => 'A') & '"'
$BIG_STRING2	'"' & (1..V-1-V/2 => 'A') & '1' & '"'
$BLANKS	(1..V-20 => ' ')
$MAX_LEN_INT_BASED_LITERAL	"2:" & (1..V-5 => '0') & "11:"
$MAX_LEN_REAL_BASED_LITERAL	"16:" & (1..V-7 => '0') & "F.E:"
$MAX_STRING_LITERAL	'"' & (1..V-2 => 'A') & '"'

```
                               1.0E308

$HIGH_PRIORITY                 7

$ILLEGAL_EXTERNAL_FILE_NAME1 \NODIRECTORY\FILENAME

$ILLEGAL_EXTERNAL_FILE_NAME2 *FLIE*

$INAPPROPRIATE_LINE_LENGTH   -1

$INAPPROPRIATE_PAGE_LENGTH   -1

$INCLUDE_PRAGMA1               PRAGMA INCLUDE ("A28006D1.TST")

$INCLUDE_PRAGMA2               PRAGMA INCLUDE ("B28006D1.TST")

$INTEGER_FIRST                -2147483648

$INTEGER_LAST                 2147483647

$INTEGER_LAST_PLUS_1          2147483648

$INTERFACE_LANGUAGE           WIN32

$LESS_THAN_DURATION           -75_000.0

$LESS_THAN_DURATION_BASE_FIRST
                              -131_073.0

$LINE_TERMINATOR              ASCII.CR & ASCII.LF

$LOW_PRIORITY                 1

$MACHINE_CODE_STATEMENT       NULL;

$MACHINE_CODE_TYPE            NO_SUCH_TYPE

$MANTISSA_DOC                 31

$MAX_DIGITS                   15

$MAX_INT                      2147483647

$MAX_INT_PLUS_1               2_147_483_648

$MIN_INT                      -2147483648

$NAME                         SHORT_SHORT_INTEGER

$NAME_LIST                    I80X86,I80386,MC680X0,S370,TRANS
                              PUTER,VAX,RS_6000,MIPS,SPARC
```

APPENDIX B

COMPILATION SYSTEM OPTIONS

The compiler options of this Ada implementation, as described in this
Appendix, are provided by the customer. Unless specifically noted
otherwise, references in this appendix are to compiler documentation and
not to this report.

```
COMPILE (SOURCE     => source_name | INSTANTIATION,
         LIBRARY    => library_name,
         OPTIONS    =>
            (ANNOTATE   => character_string,
             ERRORS     => positive_integer,
             LEVEL      => PARSE | SEMANTIC | CODE | UPDATE,
             CHECKS     => ALL | STACK | NONE,
             GENERICS   => STUBS | INLINE,
             TASKING    => YES | NO,
             MEMORY => number_of_kbytes),
         DISPLAY    =>
            (OUTPUT     => SCREEN | NONE | AUTOMATIC |
                           file_name,
             WARNING    => YES | NO,
             TEXT       => YES | NO,
             SHOW       => BANNER | RECAP | ALL | NONE,
             DETAIL     => YES | NO,
             ASSEMBLY   => CODE | MAP | ALL | NONE),
         ALLOCATION =>
            (STACK      => positive_integer),
         IMPROVE    =>
            (CALLS      => SUPPRESS | PRAGMA | AUTOMATIC,
             REDUCTION  => NONE | PARTIAL | EXTENSIVE,
             EXPRESSIONS => NONE | PARTIAL | EXTENSIVE,
         KEEP       =>
            (COPY       => YES | NO,
             DEBUG      => YES | NO,
             TREE       => YES | NO
             EDIT       => NONE | AUTOMATIC | file_name,
             OTI        => YES | NO));
```

APPENDIX C

APPENDIX F OF THE Ada STANDARD

The only allowed implementation dependencies correspond to implementation-dependent pragmas, to certain machine-dependent conventions as mentioned in Chapter 13 of the Ada Standard, and to certain allowed restrictions on representation clauses. The implementation-dependent characteristics of this Ada implementation, as described in this Appendix, are provided by the customer. Unless specifically noted otherwise, references in this Appendix are to compiler documentation and not to this report. Implementation-specific portions of the package STANDARD, which are not a part of Appendix F, are:

```
package STANDARD is

  type SHORT_SHORT_INTEGER is range -128 .. 127;
  type SHORT_INTEGER is range -32768 .. 32767;
  type INTEGER is range -2147483648 .. 2147483647;

  type FLOAT is digits 6 range
    -2#1.111_1111_1111_1111_1111_1111#E+127 ..
    2#1.111_1111_1111_1111_1111_1111#E+127

  type LONG_FLOAT is digits 15 range
-2#1.1111_1111_1111_1111_1111_1111_1111_1111_1111_1111_1111_1111_1111#E+1023
  ..
 2#1.1111_1111_1111_1111_1111_1111_1111_1111_1111_1111_1111_1111_1111#E+1023

  type DURATION is delta 2#0.000_000_000_000_01# range
    -131072.0 ..131072.0;

  end STANDARD;
```

dependent characteristics of input-output packages.

Section 9 (Characteristics of Numeric Types) Defines the ranges and attributes of numeric types in this implementation.

Section 10 (Other Implementation-Dependent Characteristics) Describes implementation-dependent characteristics not covered in the other chapters (such as that of the heap, tasks, and main subprograms).

Section 11 (Limitations) Describes compiler- and hardware-related limitations of this implementation.

Document conventions

The following list describes the typographical notations used in this document.

Italics   This font is used to designate:

               File names; for example, MAIN.CUI

               Prompts generated by a program; for example:

                  Library_Manager.NEW (LIBRARY => "\GAMES");

               (Library_Manager is the prompt.)

               Full document titles; for example, Application Developer's Guide.

               Generic command parameters in syntax diagrams (where the user must supply an actual value); for example,

                  DEFAULT.command
                  ERASE (FAMILY => family_name);

Bold   This font is used within text to designate:

               Commands that must be keyed in by the user; for example:

                  Use the command COMPILE (BINGO.ADA); to ...

Typewriter   This font is used for file listings.

The following list shows examples of actual notations used in this manual and explains how the format of the example is used to convey extra information about it.

KEEP   The underscore here indicates that KEEP is a default option.

parameter passing by the Alsys Windows NT Ada Compiler and the corresponding mechanisms of the chosen external language.

1.3 INTERFACE_NAME

Pragma INTERFACE_NAME associates the name of the interfaced subprogram with the external name of the interfaced subprogram. If pragma INTERFACE_NAME is not used, then the two names are assumed to be identical. This pragma takes the form:

    pragma INTERFACE_NAME (subprogram_name, string_literal);

where,

   subprogram_name is the name used within the Ada program to refer to the
   . interfaced subprogram.

   string_literal is the name by which the interfaced subprogram is
      referred to at link time.

The pragma INTERFACE_NAME is used to identify routines in other languages that are not named with legal Ada identifiers. Ada identifiers can only contain letters, digits, or underscores, whereas the Windows NT Linker LINK32 allows external names to contain other characters, for example, the dollar sign ($) or commercial at sign (@). These characters can be specified in the string_literal argument of the pragma INTERFACE_NAME.

The pragma INTERFACE_NAME is allowed at the same places of an Ada program as the pragma INTERFACE. (Location restrictions can be found in section 13.9 of the RM.) However, the pragma INTERFACE_NAME must always occur after the pragma INTERFACE declaration for the interfaced subprogram.

The string_literal of the pragma INTERFACE_NAME is passed through unchanged, including case sensitivity, to the COFF object file. There is no limit to the length of the name.

If INTERFACE_NAME is not used, the default link name for the subprogram is its Ada name converted to all upper case characters.

The user must be aware however, that some tools from other vendors do not fully support the standard object file format and may restrict the length or names of symbols. For example, most Windows NT debuggers only work with alphanumeric identifier names.

The Runtime Executive contains several external identifiers. All such identifiers begin with either the string "ADA_" or the string "ADAS_". Accordingly, names prefixed by "ADA_" or "ADAS_" should be avoided by the user.

Example

2.2 P'RECORD_DESCRIPTOR, P'ARRAY_DESCRIPTOR

These attributes are used to control the representation of implicit components of a record. (See Section 4.8 for more details.)

2.3 E'EXCEPTION_CODE

For a prefix E that denotes an exception name, this attribute yields a value that represents the internal code of the exception. The value of this attribute is of the type INTEGER.

2.4 Other Attributes

'OFFSET, 'RECORD_SIZE, 'VARIANT_INDEX, 'ARRAY_DESCRIPTOR, and 'RECORD_DESCRIPTOR are descibed in detail in Section 4.

Section 3

Specification of the package SYSTEM

The implementation does not allow the recompilation of package SYSTEM.

3.1 Specification of the package SYSTEM

```
--%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%
-- This unpublished work is protected both as a proprietary work and under
-- the Universal Copyright Convention and the US Copyright Act of 1976.  Its
-- distribution and access are limited only to authorized persons.  Copyright
-- (C) Alsys.  Created 1990, initially licenced 1990.  All rights reserved.
-- Unauthorized use (including use to prepare other works), disclosure,
-- reproduction or distribution may violate national criminal law.
--%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%

-- Check that all CPUs are covered.

-- Check that all operating systems are covered

package SYSTEM is

    type NAME is (I80X86,
                  I80386,
                  MC680X0,
                  S370,
                  TRANSPUTER,
                  VAX,
                  RS_6000,
                  MIPS,
```

```
--    Example:
--       "0014:00F0"
--
-- For the other targets the syntax is:
--    "00000000"   where 00000000 is an 8 digit or less hexadecimal number.
--                 For the 80386, this number represents an offset either in
--                 the data segment or in the code segment.
--                 For the MC680X0, 370 and Transputer, the number represents
--                 a virtual address (physical address for bare machines).
--    Example:
--       "00000008"
--
-- The exception CONSTRAINT_ERROR is raised if the string has not the
-- proper syntax.
-- _____


subtype ADDRESS_STRING is STRING(1..8);

function IMAGE (LEFT : in ADDRESS) return ADDRESS_STRING;
-- _____
-- Converts an address to a string.  The syntax of the returned string is
-- described in the VALUE function.
--
-- This function is used by ERROR_IO to output values of type ADDRESS.
-- Do not attempt to output an ADDRESS from within this subprogram.
-- _____


type OFFSET is range -2**31 .. 2**31 -1;
-- This type is used to measure a number of storage units (bytes). The type
-- is logically unsigned: all operations on offsets have wrap-around
-- semantics.


-- On non-segmented machines, the function and exception are meaningless.


-- _____
-- The exception CONSTRAINT_ERROR can be raised by "+" and "-".
-- _____

function "+" (LEFT : in ADDRESS; RIGHT : in OFFSET)  return ADDRESS;
function "+" (LEFT : in OFFSET;  RIGHT : in ADDRESS) return ADDRESS;
function "-" (LEFT : in ADDRESS; RIGHT : in OFFSET)  return ADDRESS;
-- _____
-- These routines provide support to perform address computations.  The
-- meaning of the "+" and "-" operators is architecture dependent.  For
-- example on a segmented machine the OFFSET parameter is added to, or
-- subtracted from the offset part of the address, the segment remaining
-- untouched.
-- _____
```

— This type is used to designate the size of an object in storage units.

```
procedure MOVE (TO      : in ADDRESS;
                FROM    : in ADDRESS;
                LENGTH  : in OBJECT_LENGTH);
```

— Copies LENGTH storage units starting at the address FROM to the address
— TO. The source and destination may overlap.
---

private

    pragma INLINE ("+", "-");

    type ADDRESS is access STRING;
    NULL_ADDRESS : constant ADDRESS := null;

end SYSTEM;

Section 4

Support for Representation Clauses

This section explains how objects are represented and allocated by the
Alsys Windows NT Ada compiler and how it is possible to control this using
representation clauses. Applicable restrictions on representation clauses
are also described.

The representation of an object is closely connected with its type. For
this reason this section addresses successively the representation of
enumeration, integer, floating point, fixed point, access, task, array and
record types. For each class of type the representation of the
corresponding objects is described.

Except in the case of array and record types, the description for each
class of type is independent of the others. To understand the
representation of array and record types it is necessary to understand
first the representation of their components.

Apart from implementation defined pragmas, Ada provides three means to
control the size of objects:

  a (predefined) pragma PACK, applicable to array types

  a record representation clause

  a size specification

For each class of types the effect of a size specification is described.
Interactions among size specifications, packing and record representation
clauses is described under the discussion of array and record types.

Size of an enumeration subtype

When no size specification is applied to an enumeration type or first named
subtype, the objects of that type or first named subtype are represented as
signed machine integers. The machine provides 8, 16 and 32 bit integers,
and the compiler selects automatically the smallest signed machine integer
which can hold each of the internal codes of the enumeration type (or
subtype). The size of the enumeration type and of any of its subtypes is
thus 8, 16 or 32 bits.

When a size specification is applied to an enumeration type, this
enumeration type and each of its subtypes has the size specified by the
length clause. The same rule applies to a first named subtype. The size
specification must of course specify a value greater than or equal to the
minimum size of the type or subtype to which it applies:

```
type EXTENDED is
    (— The usual ASCII character set.
   NUL,  SOH,  STX,  ETX,  EOT,  ENQ,  ACK,  BEL,
   ...
   'x',  'y',  'z',  '{',  '|',  '}',  '~',  DEL,

   — Extended characters
   C_CEDILLA_CAP,  U_UMLAUT,  E_ACUTE,   ...);

for EXTENDED'SIZE use 8;
— The size of type EXTENDED will be one byte. Its objects will be
— represented as unsigned 8 bit integers.
```

The Alsys compiler fully implements size specifications. Nevertheless, as
enumeration values are coded using integers, the specified length cannot be
greater than 32 bits.

Size of the objects of an enumeration subtype

Provided its size is not constrained by a record component clause or a
pragma PACK, an object of an enumeration subtype has the same size as its
subtype.

4.2 Integer Types

There are three predefined integer types in the Alsys implementation for
I80386 machines:

```
type SHORT_SHORT_INTEGER    is range -2**07 .. 2**07-1;
type SHORT_INTEGER          is range -2**15 .. 2**15-1;
type INTEGER                is range -2**31 .. 2**31-1;
```

— J is derived from SHORT_INTEGER, its size is 16 bits.

```
type N is new J range 80 .. 100;
```
— N is indirectly derived from SHORT_INTEGER, its size is
— 16 bits.

When a size specification is applied to an integer type, this integer type and each of its subtypes has the size specified by the length clause. The same rule applies to a first named subtype. The size specification must of course specify a value greater than or equal to the minimum size of the type or subtype to which it applies:

```
type S is range 80 .. 100;
for S'SIZE use 32;
```
— S is derived from SHORT_SHORT_INTEGER, but its size is
— 32 bits because of the size specification.

```
type J is range 0 .. 255;
for J'SIZE use 8;
```
— J is derived from SHORT_INTEGER, but its size is 8 bits
— because of the size specification.

```
type N is new J range 80 .. 100;
```
— N is indirectly derived from SHORT_INTEGER, but its
— size is 8 bits because N inherits the size specification
— of J.

Size of the objects of an integer subtype

Provided its size is not constrained by a record component clause or a pragma PACK, an object of an integer subtype has the same size as its subtype.

4.3 Floating Point Types

There are two predefined floating point types in the Alsys implementation for I80x86 machines:

```
type FLOAT is
    digits 6 range -(2.0 - 2.0**(-23))*2.0**127 .. (2.0 - 2.0**(-23))*2.0**127;

type LONG_FLOAT is
    digits 15 range -(2.0 - 2.0**(-52))*2.0**1023 .. (2.0 - 2.0**(-52))*2.0**1023;
```

4.3.1   Floating Point Type Representation

A floating point type declared by a declaration of the form:

```
    for FIXED'SMALL use S;

    type LONG_FIXED is delta D range (-2.0**31-1)*S .. 2.0**31*S;
    for LONG_FIXED'SMALL use S;
```

where D is any real value and S any power of two less than or equal to D.

A fixed point type declared by a declaration of the form:

```
    type T is delta D range L .. R;
```

possibly with a small specification:

```
    for T'SMALL use S;
```

is implicitly derived from a predefined fixed point type. The compiler automatically selects the predefined fixed point type whose small and delta are the same as the small and delta of T and whose range is the shortest that includes the values L to R inclusive.

In the program generated by the compiler, a safe value V of a fixed point subtype F is represented as the integer:

```
    V / F'BASE'SMALL
```

4.4.2   Fixed Point Type and Object Size

Minimum size of a fixed point subtype

The minimum possible size of a fixed point subtype is the minimum number of binary digits that is necessary for representing the values of the range of the subtype using the small of the base type.

For a static subtype, if it has a null range its minimum size is 1. Otherwise, s and S being the bounds of the subtype, if i and I are the integer representations of m and M, the smallest and the greatest model numbers of the base type such that s < m and M < S, then the minimum size L is determined as follows. For i >= 0, L is the smallest positive integer such that $I <= 2L-1$. For i < 0, L is the smallest positive integer such that $-2L-1 <= i$ and $I <= 2L-1-1$.

```
    type F is delta 2.0 range 0.0 .. 500.0;
        — The minimum size of F is 8 bits.

    subtype S is F delta 16.0 range 0.0 .. 250.0;
        — The minimum size of S is 7 bits.

    subtype D is S range X .. Y;
    — Assuming that X and Y are not static, the minimum size  of D is 7 bits
    — (the same as the minimum size of its type mark S).
```

Access Types and Objects of Access Types

The only size that can be specified for an access type using a size specification is its usual size (32 bits).

An object of an access subtype has the same size as its subtype, thus an object of an access subtype is always 32 bits long.

Collection Size

As described in RM 13.2, a specification of collection size can be provided in order to reserve storage space for the collection of an access type.

When no STORAGE_SIZE specification applies to an access type, no storage space is reserved for its collection, and the value of the attribute STORAGE_SIZE is then 0.

The maximum size is limited by the amount of memory available.

4.6 Task Types

Storage for a task activation

As described in RM 13.2, a length clause can be used to specify the storage space (that is, the stack size) for the activation of each of the tasks of a given type. Alsys also allows the task stack size, for all tasks, to be established using a Binder option. If a length clause is given for a task type, the value indicated at bind time is ignored for this task type, and the length clause is obeyed. When no length clause is used to specify the storage space to be reserved for a task activation, the storage space indicated at bind time is used for this activation.

A length clause may not be applied to a derived task type. The same storage space is reserved for the activation of a task of a derived type as for the activation of a task of the parent type.

The minimum size of a task subtype is 32 bits.

A size specification has no effect on a task type. The only size that can be specified using such a length clause is its usual size (32 bits).

An object of a task subtype has the same size as its subtype. Thus an object of a task subtype is always 32 bits long.

4.7 Array Types

Each array is allocated in a contiguous area of storage units. All the components have the same size. A gap may exist between two consecutive components (and after the last one). All the gaps have the same size.

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of its subtype:

```
type R is
   record
      K : SHORT_INTEGER;
      B : BOOLEAN;
   end record;
for R use
   record
      K at 0 range 0 .. 31;
      B at 4 range 0 .. 0;
   end record;
   -- Record type R is byte aligned. Its size is 33 bits.

type A is array (1 .. 10) of R;
   -- A gap of 7 bits is inserted after each component in order to respect
   -- the alignment of type R. The size of an array of type A will be 400 bits.
```

Array of type A: each subcomponent K has an even offset.

If a size specification applies to the subtype of the components or if the
array is packed, no gaps are inserted:


```
type R is
   record
      K : SHORT_INTEGER;
      B : BOOLEAN;
   end record;

type A is array (1 .. 10) of R;
pragma PACK(A);
   -- There is no gap in an array of type A because A is packed.
   -- The size of an object of type A will be 330 bits.

type NR is new R;
for NR'SIZE use 24;

type B is array (1 .. 10) of NR;
   -- There is no gap in an array of type B because
   -- NR has a size specification.
   -- The size of an object of type B will be 240 bits.
```


4.7.2   Array Subtype and Object Size

Size of an array subtype

The size of an array subtype is obtained by multiplying the number of its
components by the sum of the size of the components and the size of the
gaps (if any). If the subtype is unconstrained, the maximum number of
components is considered.

A record representation clause need not specify the position and the size for every component. If no component clause applies to a component of a record, its size is the size of its subtype.

4.8.2   Indirect Components

If the offset of a component cannot be computed at compile time, this offset is stored in the record objects at run time and used to access the component. Such a component is said to be indirect while other components are said to be direct.

If a record component is a record or an array, the size of its subtype may be evaluated at run time and may even depend on the discriminants of the record. We will call these components dynamic components:

```
type DEVICE is (SCREEN, PRINTER);

type COLOR is (GREEN, RED, BLUE);

type SERIES is array (POSITIVE range <>) of INTEGER;

type GRAPH (L : NATURAL) is
    record
        X : SERIES(1 .. L); -- The size of X depends on L
        Y : SERIES(1 .. L); -- The size of Y depends on L
    end record;

Q : POSITIVE;

type PICTURE (N : NATURAL; D : DEVICE) is
    record
        F : GRAPH(N); -- The size of F depends on N
        S : GRAPH(Q); -- The size of S depends on Q
        case D is
            when SCREEN =>
                C : COLOR;
            when PRINTER =>
                null;
        end case;
    end record;
```

Any component placed after a dynamic component has an offset which cannot be evaluated at compile time and is thus indirect. In order to minimize the number of indirect components, the compiler groups the dynamic components together and places them at the end of the record:

The record type PICTURE: F and S are placed at the end of the record

Note that Ada does not allow representation clauses for record components with non-static bounds [RM 13.4.7], so the compiler's grouping of dynamic components does not conflict with the use of representation clauses.

(note that the storage effectively allocated for the record object may be more than this).

The value of a RECORD_SIZE component may denote a number of bits or a number of storage units. In general it denotes a number of storage units, but if any component clause specifies that a component of the record type has an offset or a size which cannot be expressed using storage units, then the value designates a number of bits.

The implicit component RECORD_SIZE must be large enough to store the maximum size of any value of the record type. The compiler evaluates an upper bound MS of this size and then considers the implicit component as having an anonymous integer type whose range is 0 .. MS.

If R is the name of the record type, this implicit component can be denoted in a component clause by the implementation generated name R'RECORD_SIZE. This allows user control over the position of the implicit component in the record.

VARIANT_INDEX

This implicit component is created by the compiler when the record type has a variant part. It indicates the set of components that are present in a record value. It is used when a discriminant check is to be done.

Component lists in variant parts that themselves do not contain a variant part are numbered. These numbers are the possible values of the implicit component VARIANT_INDEX.

```
type VEHICLE is (AIRCRAFT, ROCKET, BOAT, CAR);

type DESCRIPTION (KIND : VEHICLE := CAR) is
   record
      SPEED : INTEGER;
      case KIND is
         when AIRCRAFT | CAR =>
             WHEELS : INTEGER;
             case KIND is
                when AIRCRAFT =>        -- 1
                    WINGSPAN : INTEGER;
                when others =>  -- 2
                    null;
             end case;
         when BOAT => -- 3
            STEAM  : BOOLEAN;
         when ROCKET =>        -- 4
            STAGES : INTEGER;
      end case;
   end record;
```

The value of the variant index indicates the set of components that are present in a record value.

effect. The only size that can be specified using such a length clause is its usual size. Nevertheless, such a length clause can be useful to verify that the layout of a record is as expected by the application.

Size of an object of a record subtype

An object of a constrained record subtype has the same size as its subtype.

An object of an unconstrained record subtype has the same size as its subtype if this size is less than or equal to 8 kb. If the size of the subtype is greater than this, the object has the size necessary to store its current value; storage space is allocated and released as the discriminants of the record change.

Section 5

Conventions for Implementation-Generated Names

The Alsys Windows NT Ada Compiler may add fields to record objects and have descriptors in memory for record or array objects. These fields are accessible to the user through implementation-generated attributes (See Section 2.3).

The following predefined packages are reserved to Alsys and cannot be recompiled in Version 5.5:

    system
    calendar
    internal_types
    system_environment
    interrupt_manager
    unix_types
    unsigned
    machine_operations_386
    get_file_number
    alsys_codegen_support
    alsys_rts_extended_ascii
    alsys_traces
    alsys_target_integers
    alsys_rt_types
    alsys_time_types
    alsys_machine_task_types
    alsys_stack_extension
    alsys_tcb_package
    alsys_assert
    alsys_task_lists
    alsys_resource
    alsys_synchronization
    alsys_ada_runtime
    alsys_error_io
    alsys_machine

6.2 Address Clauses for Program Units

Address clauses for program units are not implemented in the current version of the compiler.

6.3 Address Clauses for Interrupt Entries

Interrupt entries are not supported.

Section 7

Unchecked Conversions

Unchecked conversions are allowed between any types provided the instantiation of UNCHECKED_CONVERSION is legal Ada. It is the programmer's responsibility to determine if the desired effect is achieved.

If the target type has a smaller size than the source type then the target is made of the least significant bits of the source.

Section 8

Input-Output Packages

In this part of the Appendix the implementation-specific aspects of the input-output system are described.

8.1 Introduction

In Ada, input-output operations (IO) are considered to be performed on objects of a certain file type rather than being performed directly on external files. An external file is anything external to the program that can produce a value to be read or receive a value to be written. Values transferred for a given file must be all of one type.

Generally, in Ada documentation, the term file refers to an object of a certain file type, whereas a physical manifestation is known as an external file. An external file is characterized by

Its name, which is a string defining a legal path name under the current version of the operating system.

Its form, which gives implementation-dependent information on file characteristics.

Both the name and the form appear explicitly as parameters of the Ada CREATE and OPEN procedures. Though a file is an object of a certain file

| COUNT | 0 .. 2147483647 | — $2**31 - 1$ |
| POSITIVE_COUNT | 1 .. 2147483647 | — $2**31 - 1$ |

For the package TEXT_IO, the range of values for the type FIELD is as follows:

| FIELD | 0 .. 255 | — $2**8 - 1$ |

9.2 Floating Point Type Attributes

	FLOAT	LONG_FLOAT
DIGITS	6	15
MANTISSA	21	51
EMAX	84	204
EPSILON	9.53674E-07	8.88178E-16
LARGE	1.93428E+25	2.57110E+61
SAFE_EMAX	125	1021
SAFE_SMALL	1.17549E-38	2.22507E-308
SAFE_LARGE	4.25353E+37	2.24712E+307
FIRST	-3.40282E+38	-1.79769E+308
LAST	3.40282E+38	1.79769E+308
MACHINE_RADIX	2	2
MACHINE_EMAX	128	1024
MACHINE_EMIN	-125	-1021
MACHINE_ROUNDS	true	true
MACHINE_OVERFLOWS	false	false
SIZE	32	64

9.3 Attributes of Type DURATION

| DURATION'DELTA | $2.0 ** (-14)$ |
| DURATION'SMALL | $2.0 ** (-14)$ |

memory (hard disk swap space).

10.3 Characteristics of Tasks

The default task stack size is 4K bytes (96K bytes for the environment task), but by using the Binder option STACK.TASK the size for all task stacks in a program may be set to a size from 1K bytes to 32767 bytes.

Preemption of Ada tasks are performed by Windows NT since they are Windows NT threads. PRIORITY values are in the range 1..86. A task with undefined priority (no pragma PRIORITY) will take the default priority given by Windows NT.

The accepter of a rendezvous executes the accept body code in its own stack. Rendezvous with an empty accept body (for synchronization) does not cause a context switch.

The main program waits for completion of all tasks dependent upon library packages before terminating.

Abnormal completion of an aborted task takes place immediately, except when the abnormal task is the caller of an entry that is engaged in a rendezvous, or if it is in the process of activating some tasks. Any such task becomes abnormally completed as soon as the state in question is exited.

The message

Deadlock in Ada program

is printed to STANDARD_ERROR when the Runtime Executive detects that no further progress is possible for any task in the program. The execution of the program is then abandoned.

10.4 Definition of a Main Subprogram

A library unit can be used as a main subprogram if and only if it is a procedure that is not generic and that has no formal parameters.

10.5 Ordering of Compilation Units

The Alsys Windows NT Ada Compiler imposes no additional ordering constraints on compilations beyond those required by the language.

Section 11

Limitations

11.1 Compiler Limitations